



Bug classification using r2

beebug: A tool for checking exploitability

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- ☐ I have never found any web vulnerability 😊

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Disclaimer: All the work presented here is mine (not of my employer)

- Motivation
- Common vulnerabilities
- Exploitable crashes
- Tool and usage
- Graph generation
- Demo
- Future direction
- Thank you

Motivation

- If you work as vulnerability researcher, often you have to analyze a crash in order to understand:
 - what type of vulnerability you have found
 - whether it is exploitable
- How to **quickly** determine if a crash in a target is potentially exploitable
- Existing similar tools:
 - *exploitable.py* for gdb
 - *!exploitable* for WinDbg
- However, the tool presented here is easy to integrate into a automated script

Common Vulnerabilities

- Heap Overflow (e.g., CVE 2018-11625)
- UAF (e.g., CVE-2018-11624)
- Stack Overflow (e.g., CVE-2018-12327)
- Integer Overflow (e.g., CVE-2017-15587)
- Type confusion (e.g., CVE-2018-4944)
- Out-of-bounds (e.g., CVE-2018-6149)

- Stack Overflow - **CVE-2018-12327** (ntpq 4.2.8p11 Local BufferOverflow)

```
root@invictus1306-VirtualBox:/home/invictus1306/Downloads/ntp-4.2.8p11# ./ntpd/ntpd -4 ['python -c 'print "A" * 300']
Name or service not known
*** stack smashing detected ***: ./ntpd/ntpd terminated
Aborted (core dumped)
```

```
#define LENHOSTNAME 256 /* host name is 256 characters long */
```

```
static int
openhost(
    const char *hname,
    int fam
)
{
    const char svc[] = "ntp";
    char temphost[LENHOSTNAME];
    int a_info, i;
    struct addrinfo hints, *ai;
    sockaddr_u addr;
    size_t octets;
    register const char *cp;
    char name[LENHOSTNAME];

    /*
     * We need to get by the [] if they were entered
     */

    cp = hname;

    if (*cp == '[') {
        cp++;
        for (i = 0; *cp && *cp != ']'; cp++, i++)
            name[i] = *cp;
        if (*cp == ']') {
            name[i] = '\0';
        }
    }
}
```



Crash!!

- Some crashes are not easy to classify and many times these are classified as Null Pointer Dereference that lead to “Denial of Service”
- A Null Pointer Dereference vulnerability could also potentially lead to code execution.

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SAPCAR Heap Buffer Overflow: From crash to exploit

BY: Maximiliano Vidal  Latest from CoreLabs

1. Introduction



In this blog post, we will cover the analysis and exploitation of a [simple heap buffer overflow](#) found in SAPCAR a few weeks ago.

SAP published security note #2441560 classifying the issue as "Potential Denial of Service". This post is our attempt to ~~show that code execution is not only possible but also relatively easy to achieve. The idea is to provide a (hopefully!) cohesive example for other beginners that might be interested in binary exploitation. We will see one possible approach to make sense out of a few hundred crashes obtained through *fuzzing*, how to identify the root cause of the bug, and how to determine its exploitability. Afterwards, we will develop an exploit using the old and well known file pointer overwrite technique. The last section will go into some more detail about a relevant mitigation implemented in *glibc 2.24*.~~

- Sometime a crash report is rejected because it seems to be not exploitable



FROM CRASH TO EXPLOIT: CVE-2015-6086 – OUT OF BOUND READ/ASLR BYPASS

18/01/2016 / 5 Comments on From Crash to Exploit: CVE-2015-6086 – Out of Bound Read/ASLR Bypass / in Research / by admin_payatu

INTRODUCTION

This is a story of an **Out of Bound Read** bug in **Internet Explorer 9-11**. This is almost **5 years old bug** which got discovered in **April 2015**. It is a very interesting bug, at least from my perspective, because it was **rejected** almost **4-5 times** by Zero Day Initiative (ZDI) stating that it's **not exploitable**.

My **SVG fuzzer** was hitting a crash continuously, at first, the bug looked like usual **Use after Free** as it was trying to **read Invalid Memory**. But after triaging it turned out to be **Out of Bound Read** bug. I submitted this bug to **ZDI** and they rejected it at first stating that they are not able to reproduce it. Later they rejected

- For example if a vulnerability is reported as Denial of Service, the patch may not provide the proper mitigation
- The same vulnerability might be able to give code execution if analyzed/exploited properly

Exploitable crashes

Some crashes that are generally exploitable:

- Stack overflow
- Crash on Program Counter
- Crash on branch
- Crash on write memory
- Heap vulnerabilities
- Read access violation ??

Stack overflow (1/2)

When a stack buffer overflow occurs – e.g., overwriting control information such as a return address - libc (depending on the *gcc* version) generates a signal and stops the target:

```
*** stack smashing detected ***: ./stack_overflow terminated
Aborted (core dumped)
```

Generally this type of stack buffer overflow is considered exploitable – e.g., the size of the buffer is fixed and there is weak bounds checking

How is it possible to detect?

`__stack_chk_fail` terminates a function in case of stack overflow:

```
[0x7fea3e5ee428]> pd -1 @ sym.exploit_me+123
0x0040060c e85ffeffff call sym.imp.__stack_chk_fail
```

Stack overflow (2/2)

Look at the backtrace:

```
==13765==  
*** stack smashing detected ***: ./stack_overflow terminated  
==13765==  
==13765== Process terminating with default action of signal 6 (SIGABRT)  
==13765==   at 0x4E6F428: raise (raise.c:54)  
==13765==   by 0x4E71029: abort (abort.c:89)  
==13765==   by 0x4EB17E9: __libc_message (libc_fatal.c:175)  
==13765==   by 0x4F5315B: __fortify_fail (fortify_fail.c:37)  
==13765==   by 0x4F530FF: __stack_chk_fail (stack_chk_fail.c:28)  
==13765==   by 0x400010: exploit_me (in /home/invictus1306/Documents/r2conf/init_test/mytests/stack_overflow)  
==13765==   by 0x400040: ??? (in /home/invictus1306/Documents/r2conf/init_test/mytests/stack_overflow)  
==13765==   by 0xFFFF0003F7: ???  
==13765==   by 0xFFFFFFFF: ???  
==13765==   by 0x40063F: ??? (in /home/invictus1306/Documents/r2conf/init_test/mytests/stack_overflow)  
==13765==   by 0x4E5A82F: (below main) (libc-start.c:291)  
==13765==
```

Looking at the backtrace it is possible to understand that the target program called the function `__stack_chk_fail`, which terminated it with the **SIGABRT** signal.

Another thing that is possible to see in the backtrace is the presence of the `__fortify_fail` function, it means that the StackGuard patch to GCC is there.

gcc version 5.4.0 20160609 (Ubuntu 5.4.0-6ubuntu1~16.04.9)

Crash on Program Counter (1/2)

If there is a **SIGSEGV** signal and the address of the program counter is equal to the address of the memory error, this could tell us that the *PC* is tainted and can be controlled by an attacker.

Generally this type of crash is considered exploitable – e.g., the presence of a call instruction with a bad argument.

In the following case we have a **SIGSEGV** signal and we can see that the address of the “memory error” is equal to the “program counter” address.

```
[0x7f10e7b52c30]> dc  
child stopped with signal 11  
[+] SIGNAL 11 errno=0 addr=0x00601038 code=2 ret=0  
[0x00601038]> █
```

```
[0x00601038]> dr rip  
0x00601038  
[0x00601038]> █
```

Let's go deeper...

Crash on Program Counter (2/2)

Let's see the content of the current instruction:

```
Search your computer
[0x00001038]> pd -1
;-- foo:
;-- rax:
;-- rip:
0x00601038 0000 add byte [rax], al
```

Uhhh *NULL* bytes, but who is the caller?

```
[0x00601038]> dbt
0 0x601038 sp: 0x0 0 [??] obj.foo obj.foo0
1 0x4004f1 sp: 0x7ffd08381fa8 0 [sym.main] main+27
2 0x7f10e77a8830 sp: 0x7ffd08381fc8 32 [??] r11+240
3 0x7f10e7b627cb sp: 0x7ffd08382058 144 [??] sym.dl_rtl_d_i_serinfo+29051
4 0x400409 sp: 0x7ffd08382088 48 [??] entry0+41
[0x00601038]> pd -1 @ 0x4004f1
| 0x004004ef ffd0 call rax
[0x00601038]>
```

As we suspected there is a call statement, which has *rax* as parameter, this could mean that the *rax* value could be controlled by the attacker and as consequence it could be exploitable.

Crash on Branch

A crash on a branch instruction could indicate that the control flow could be controlled by the attacker .

For example, in that case:

```
[0x004004e1]> dc
child stopped with signal 11
[+] SIGNAL 11 errno=0 addr=0xffffffff code=1 ret=0
[0x004004e1]> pd 1
;-- rip:
0x004004e1 ff1425ffffff. call qword [0xffffffff]
[0x004004e1]> █
```

We have a **SIGSEGV** signal, and the current instruction is a call to an address that is not mapped in the memory.

Crash on write memory

An access violation on a destination operand could indicate a write access violation.

```
2000
[0x7f2f01f32c30]> dc
child stopped with signal 11
[+] SIGNAL 11 errno=0 addr=0x00010000 code=1 ret=0
[0x004004eb]> pd 1
;-- rip:
0x004004eb c60041 mov byte [rax], 0x41 ; 'A' ; [0x41:1]=255 ; 65
[0x004004eb]> dc
```

Generally this type of crash is considered exploitable – e.g., an attacker may control the write address/value.

Note: The faulting address is equal to the destination operand value:

Faulting address: **0x10000**

Fault instruction: **mov byte [rax], 0x41**

rax = **0x10000**

Heap vulnerabilities

The newer version of *glibc* added a lot of sanity checks for chunk headers:

```
[#0] 0x7ffff7a42428 → Name: __GI_raise(sig=0x6)
[#1] 0x7ffff7a4402a → Name: __GI_abort()
[#2] 0x7ffff7a847ea → Name: __libc_message(do_abort=0x2, fmt=0x7ffff7b9ded8 "*** Error in '%s': %s: 0x%s ***\n")
[#3] 0x7ffff7a8d37a → Name: malloc_printerr(ar_ptr=<optimized out>, ptr=<optimized out>, str=0x7ffff7b9df50 'free(): invalid next size (fast)', action=0x3)
[#4] 0x7ffff7a8d37a → Name: __int_free(av=<optimized out>, p=<optimized out>, have_lock=0x0)
[#5] 0x7ffff7a9153c → Name: __GI___libc_free(mem=<optimized out>)
[#6] 0x4005f5 → Name: main()
```

```
[#0] 0x7ffff7a42428 → Name: __GI_raise(sig=0x6)
[#1] 0x7ffff7a4402a → Name: __GI_abort()
[#2] 0x7ffff7a847ea → Name: __libc_message(do_abort=0x2, fmt=0x7ffff7b9ded8 "*** Error in '%s': %s: 0x%s ***\n")
[#3] 0x7ffff7a8d37a → Name: malloc_printerr(ar_ptr=<optimized out>, ptr=<optimized out>, str=0x7ffff7b9dfa0 "double free or corruption (fasttop)", act
[#4] 0x7ffff7a8d37a → Name: __int_free(av=<optimized out>, p=<optimized out>, have_lock=0x0)
[#5] 0x7ffff7a9153c → Name: __GI___libc_free(mem=<optimized out>)
[#6] 0x400762 → Name: main()
```

Looking at the backtrace it is possible to see the target program calls the function `__int_free`, and also see the `malloc_printerr` function contains the descriptive string.

This helps the analyzer to understand the type of vulnerability.

Read access violation that leads to code execution

As mentioned, sometimes a simple read access violation vulnerability could lead to code execution.

```
0x401bfd <main+1177>    mov     rdi, rax
0x401c00 <main+1180>    call   0x401dbc <_ZN4Base3setEff>
0x401c05 <main+1185>    mov     rax, QWORD PTR [rbp-0x68]
0x401c09 <main+1189>    mov     rax, QWORD PTR [rax]
0x401c0c <main+1192>    add     rax, 0x8
→ 0x401c10 <main+1196>    mov     rax, QWORD PTR [rax]
0x401c13 <main+1199>    mov     rdx, QWORD PTR [rbp-0x08]
0x401c17 <main+1203>    mov     rdi, rdx
0x401c1a <main+1206>    call   rax
0x401c1c <main+1208>    movd    eax, xmm0
0x401c20 <main+1212>    mov     DWORD PTR [rbp-0x84], eax
```

```
Program received signal SIGSEGV, Segmentation fault.
0x0000000000401c10 in main ()
```

```
$rax : 0x0000000000000008
$rbx : 0x0000000000000000 → 0x0000000000000000
```

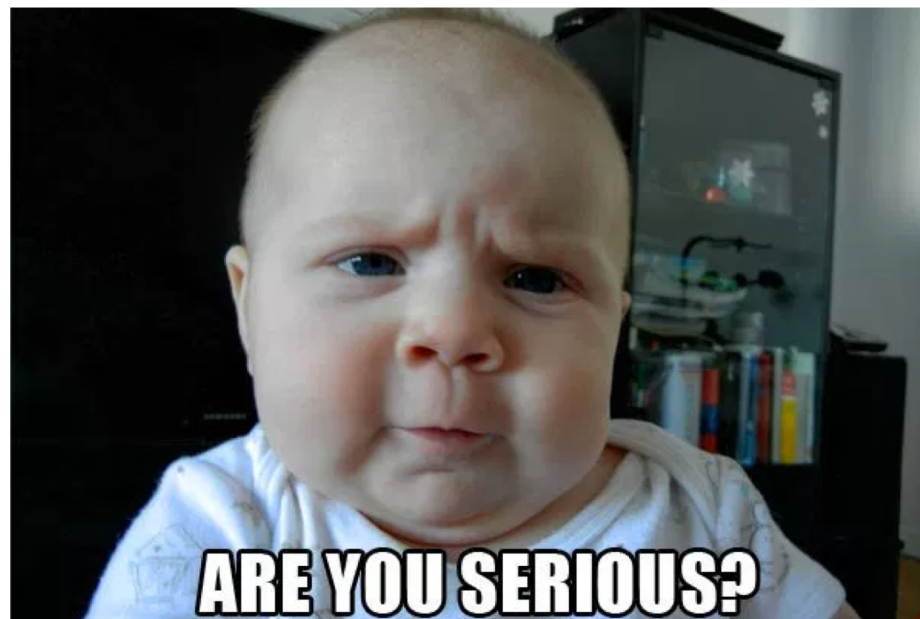
The “exploitable” gdb plugin tell us the crash is `PROBABLY_NOT_EXPLOITABLE`

```
gef> exploitable
__main__:99: UserWarning: GDB v7.11 may not support required Python API
Description: Access violation near NULL on source operand
Short description: SourceAvNearNull (16/22)
Hash: c0f280e80c376bd06c228bee621e541.c0f280e80c376bd06c228bee621e541
Exploitability Classification: PROBABLY_NOT_EXPLOITABLE
Explanation: The target crashed on an access violation at an address matching the source operand of the current instruction. This likely indicates a read access violation, which may mean the application crashed on a simple NULL dereference to data structure that has no immediate effect on control of the processor.
Other tags: AccessViolation (21/22)
```


Read access violation that leads to code execution

Instead this is exploitable, in this case the *rax* register is controllable from an attacker.

```
[0x00401c00]> s 0x401c10
[0x00401c10]> pd 10
0x00401c10      488b00      mov rax, qword [rax]
0x00401c13      488b5598    mov rdx, qword [rbp - 0x68]
0x00401c17      4889d7      mov rdi, rdx
0x00401c1a      ffd0      call rax
0x00401c1c      660f7ec0    movd eax, xmm0
```



Tool and usage

New tool released as open source:

beebug

<https://github.com/invictus1306/beebug>

Some features:

- Classify if a crash could be exploitable:
 - Stack overflow on libc
 - Crash on Program Counter
 - Crash on branch
 - Crash on write memory
 - Heap vulnerabilities
 - Read access violation (some exploitable cases)
- Help to analyze a crash (graph view)

Help view

```
root@invictus1306-VirtualBox:/home/invictus1306/Documents/r2conf/beebug/beebug# python3 beebug.py -h
usage: beebug.py [-h] -t TARGET [-a TARGETARGS] [-f FILE] [-g GRAPH]
```

optional arguments:

```
-h, --help            show this help message and exit
-t TARGET, --target TARGET
                        target program to analyze
-a TARGETARGS, --targetargs TARGETARGS
                        arguments for the target program
-f FILE, --file FILE  input file
-g GRAPH, --graph GRAPH
                        generate the graph
```

Output example

Crash on PC - Generally it is exploitable, the PC could be tainted

backtrace

0	0x601038	sp: 0x0	0	[??]	obj.foo obj.foo0
1	0x4004f1	sp: 0x7fff447da688	0	[sym.main]	main+27
2	0x7fba9cd9f830	sp: 0x7fff447da6a8	32	[??]	r11+240
3	0x7fba9d1597cb	sp: 0x7fff447da738	144	[??]	sym.dl_rtl_d_serinfo+29051
4	0x400409	sp: 0x7fff447da768	48	[??]	entry0+41

registers

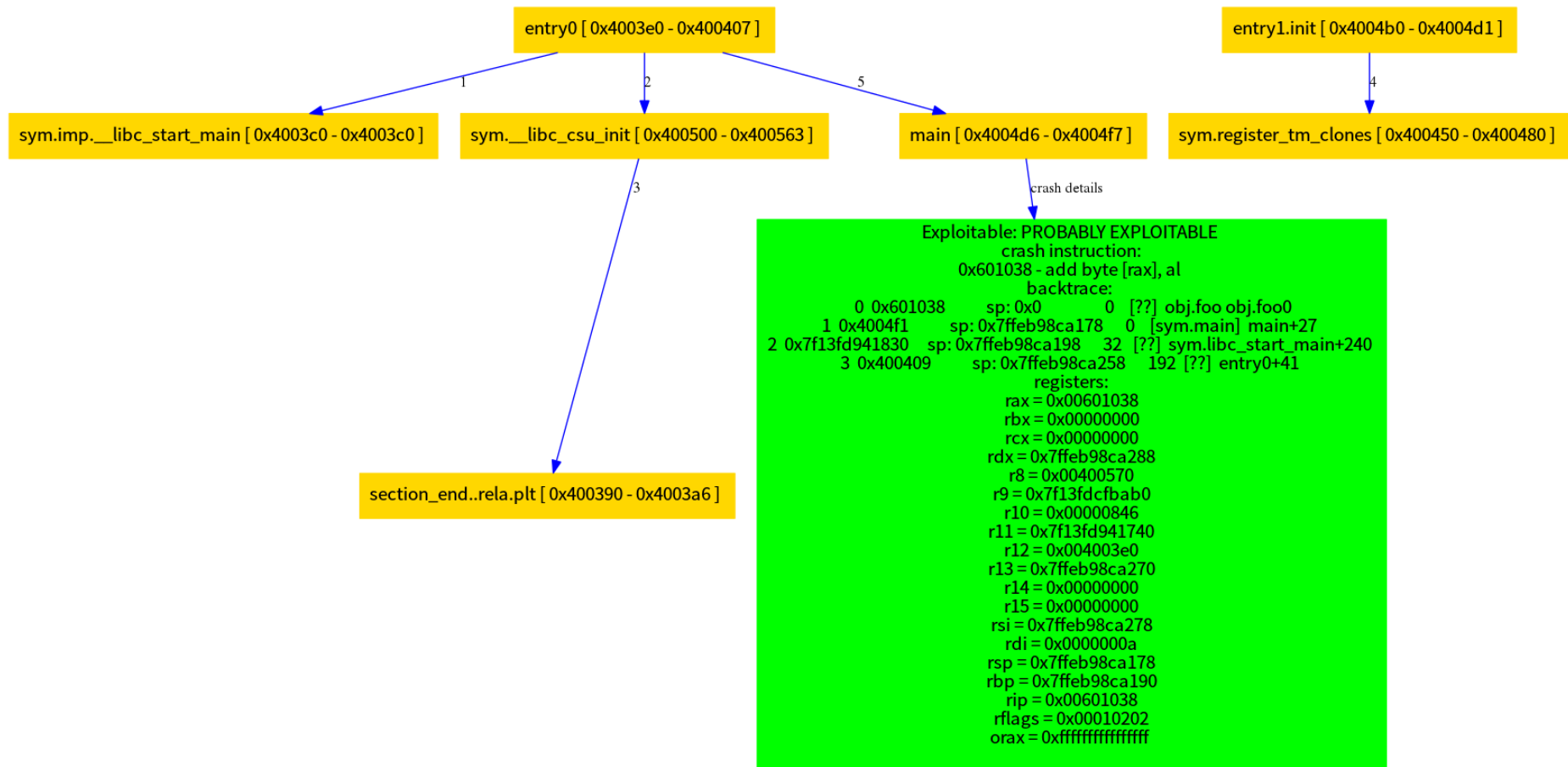
```
rax = 0x00601038
rbx = 0x00000000
rcx = 0x00000000
rdx = 0x7fff447da798
r8 = 0x00400570
r9 = 0x7fba9d159ab0
r10 = 0x00000846
r11 = 0x7fba9cd9f740
r12 = 0x004003e0
r13 = 0x7fff447da780
r14 = 0x00000000
r15 = 0x00000000
```

- ```
root@invictus1306-VirtualBox:/home/invictus1306/Documents/r2conf/beebug# python3 beebug.py -t /home/invictus1306/Downloads/ntp-4.2.8p11/ntpd/ntpd -a "-4 [`python -c 'print \"A\" * 300`']"
```
- ```
Process with PID 5017 started...  
File dbg:///home/invictus1306/Downloads/ntp-4.2.8p11/ntpd/ntpd -4 [AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
AAA] reopened in read-write mode  
= attach 5017 5017  
WARNING: bin_strings buffer is too big (0xffffaa1e5556a5f8). Use -zzz or set bin.maxstrbuf (RABIN2_MAXSTRBUF) in r2 (rabin  
2)  
WARNING: bin_strings buffer is too big (0xffffaa1e5555a258). Use -zzz or set bin.maxstrbuf (RABIN2_MAXSTRBUF) in r2 (rabin  
2)  
Cannot find function 'entry0' at 0x00007bd0  
Name or service not known  
*** stack smashing detected ***: /home/invictus1306/Downloads/ntp-4.2.8p11/ntpd/ntpd terminated  
child stopped with signal 6  
[+] SIGNAL 6 errno=0 addr=0x00001399 code=-6 ret=0
```
- ```
Stack error - Generally it could be exploitable, due to suspicious functions in the backtrace
```
- ```
backtrace  
0 0x7fc8112a8428    sp: 0x0                0   [??]  
1 0x7fc8112aa02a    sp: 0x7ffcbbfab54e8    0   [??] sym.abort+362  
2 0x7fc8112ea7ea    sp: 0x7ffcbbfab5618    304 [??] sym.fsetlocking+778  
3 0x7fc8112ea5ce    sp: 0x7ffcbbfab56a8    144 [??] sym.fsetlocking+238
```

Graph generation

- As already mentioned, some crashes are not easy to classify
- With **beebug** there is the possibility to have a runtime graph view
- The sequence of the executed functions will be printed out
- This could help the analyzer to have a better understanding of what is happening
- For every node in the graph this is the label:
 - Function name [start_address – end_address]
 - The edge, contains the number in the sequence of execution

Graph generation



DEMO

Future direction

- Support different architectures
- Improvement of the graph view (based on radare2)
- Analyze core dumps (based on radare2)
- Use instrumentation for the graph view generation

beebug

<https://github.com/invictus1306/beebug>

Thank you!

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